An Online Learning Approach to Improving the Quality of Crowd-Sourcing

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The power of crowdsourcing

Data collection: participatory sensing, user-generated map :



Recommendation: rating of movies, news, restaurants, services:



Social studies: opinion survey, the science of opinion survey:



Crowd-sourcing market

Data processing: image labeling, annotation



- Paid workers perform computational tasks.
- Hard to measure and evaluate quality objective: competence, bias, irresponsible behavior, etc.

One step further: Use the wisdom of crowd



- Redundant assignment.
- Label aggregation.

Our vision:



- Labeler selection. (first step)
- Adaptive learning. (second step)
- A weighted version of selection. (one more step)

Our objective

To make the most effective use of the crowdsourcing system

- Cost in having large amount of data labeled is non-trivial
- Time constrained machine learning tasks.

A sequential/online learning framework

- Over time learn which labelers are more competent, or whose reviews/opinion should be valued more.
- Quality control rather than random assignment
- Closed-loop, causal.

Multiarmed bandit (MAB) framework

A sequential decision making and learning framework:

- Objective: select the best of a set of choices ("arms")
- Principle: repeated sampling of different choices ("exploration"), while controlling how often each choice is used based on their empirical quality ("exploitation").
- Performance measure: "regret" difference between an algorithm and a benchmark.



Challenges and key ideas

Main challenge in crowdsourcing: ground truth

- True label of data remains unknown
- If view each labeler as a choice/arm: unknown quality of outcome ("reward").

Key features:

- Mild assumption on the collective quality of the crowd; quality of an individual is estimated against the crowd.
- Online learning: Learning occurs as data/labeling tasks arrive.
- Comparing against optimal static selections.

Outline of the talk

Problem formulation

Online solution

Simple/weighted majority voting

Extensions and discussions

Experiment results

Numerical results & Results on AMT data

Conclusion and on-going works

Labeler selection

M labelers; labelers *i* has accuracy p_i (can be task-dependent).

- ▶ No two exactly the same: $p_i \neq p_j$ for $i \neq j$, and $0 < p_i < 1$, $\forall i$.
- Collective quality: $\bar{p} := \sum_i p_i / M > 1/2$.
- ▶ Probability that a simple majority vote over all M labelers is correct: a_{min} := P(∑_i X_i/M > 1/2).

• If
$$\bar{p} > 1/2$$
 and $M > \frac{\log 2}{\bar{p} - 1/2}$, then $a_{\min} > 1/2$.

Unlabeled tasks arrive at $t = 1, 2, \cdots$.

- User selects a subset S_t of labelers for task at t.
- Labeling payment of c_i for each task performed by labeler *i*.

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Labeling outcome/Information aggregation



Aggregating results from multiple labelers:

- A task receives a set of labels: $\{L_i(t)\}_{i \in S_t}$.
- Use simple majority voting to compute the label output: L*(t). (extensible to weighted majority voting)

Probability of correct labeling outcome: $\pi(S_t)$.

• Optimal set of labelers: S^* that maximizes $\pi(S)$.

$$\pi(S_t) = \underbrace{\sum_{\substack{S:S \subseteq S_t, |S| \ge \lceil \frac{|S_t|+1}{2} \rceil}} \prod_{i \in S} p_i \cdot \prod_{j \in S_t \setminus S} (1-p_j)}_{\text{Majority wins}} + \underbrace{\frac{\sum_{S:S \subseteq S_t, |S| = \frac{|S_t|}{2}} \prod_{i \in S} p_i \cdot \prod_{j \in S_t \setminus S} (1-p_j)}{2}_{\text{Ties broken equally likely}}}$$

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Assuming known $\{p_i\}$, S^* can be obtained using a linear search

Theorem

Under the simple majority voting rule, $|S^*|$ is an odd number. Furthermore, S^* is monotonic: if $i \in S^*$ and $j \notin S^*$, then we must have $p_i > p_j$.

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Objective of our online solution

Labeler expertise p_i s being unknown a priori

 Goal: Gradually learn labelers' quality and make selections adaptively

Performance measure:

Comparing with the optimal selection (static):

$$R(T) = T\pi(S^*) - E[\sum_{t=1}^T \pi(S_t)]$$

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An online learning algorithm



Interlveaing explorations of MAB and Crowd-sourcing: Double exploration

- There is a set of tasks E(t) (~ log t) used for *testing* purposes.
- ► These or their independent and identical variants are repeatedly assigned to the labelers (~ log t).¹

¹More discussions follow later on independence.





Two types of time steps:

- ▶ Double exploration: all M labelers are used. Exploration is entered if (1) the number of testers falls below a threshold (~ log t), or if (2) the number of times a tester has been tested falls below a threshold (~ log t).
- Exploitation: the estimated S
 ^{*} is used to label the arriving task based on the current estimated {p
 _i}.





Three types of tasks:

- ► Testers: those arriving to find (1) true and (2) false. These are added to E(t) and are repeatedly used to collect independent labels whenever (2) is true subsequently.
- Throw-aways: those arriving to find (2) true. These are given a random label.
- ▶ Keepers: those arriving to find both (1) and (2) false. These are given a label outcome using the best estimated set of labelers.





Accuracy update

- Estimated label on tester k at time t: majority label over all test outcomes up to time t.
- *p˜_i* at time *t*: the % of times *i*'s label matches the majority vote known at *t* out of all tests on all testers.

Regret

Main result:

$$\mathsf{R}(\mathsf{T}) \hspace{.1in} \leq \hspace{.1in} \mathsf{Const}(\mathsf{S}^*,\Delta_{\mathsf{max}},\Delta_{\mathsf{min}},\delta_{\mathsf{max}},\delta_{\mathsf{min}},\mathsf{a_{\mathsf{min}}})\log^2(\mathsf{T}) + \mathsf{Const}$$

$$\Delta_{\max} = \max_{S \neq S^*} \pi(S^*) - \pi(S), \ \delta_{\max} = \max_{i \neq j} |p_i - p_j|.$$

- First term due to exploration; second due to exploitation.
- Can obtain similar result on the cost C(T).

Weighted majority voting



• Each labeler *i*'s decision is weighed by $\log \frac{p_i}{1-p_i}$.

Theorem

Under the weighted majority vote and assuming $p_i \ge 0.5, \forall i$, the optimal set S^* is monotonic, i.e., if we have $i \in S^*$ and $j \notin S^*$ then we must have $p_i > p_j$.

Main results on weighted majority voting:

- $R(T) \leq O(\log^2 T)$, but with strictly larger constants.
- Have to account for additional error in estimating the weights when determining label outcome.
- ► A larger constant: slower convergence to a better target.

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Re-assignment of the testers

IID noise insertion



Random delay

- Commonly adopted in survey methodology to ensure valid responses.
- ▶ Bounded delay τ_{max} leads to τ_{max} -fold regret.

Other extensions

Prior knowledge on several constants

- Exploration length depends on several system parameters.
- Sub-logarithmic remedy without knowing prior knowledge.

Improve the bound by improving a_{\min} : weed out bad labelers.

- Ranking based on counting of disagreement.
- Start the weeding out from the end of list.
- Requires only $O(\log T)$ samples to achieve a bounded regret.

Other extensions cont.

Labelers with different type of tasks

- ► Finite number of types: Similar results.
- Infinite number of types.
 - Continuous MAB.
 - Sub-linear regret bound.

With delayed arrival of ground-truth

► O(log T) regret time uniformly.

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Experiment I: simulation with M = 5





Effect of a_{\min} : higher a_{\min} leads to much better performance.



Performance comparison

labeler selection v.s. full crowd-sourcing (simple majority vote)



Comparing weighted and simple majority vote



Table: Average reward per labeler: there is a clear gap between with and without using LS_OL.

Experiment II: on a real AMT dataset

The dataset

- Contains 1,000 images each labeled by the same set of 5 AMTs.
- Labels are on a scale from 0 to 5, indicating how many scenes are seen from each image.
- A second dataset summarizing keywords for scenes of each image: use this count as the ground truth.

Counting number of disagreement (online):

	AMT1	AMT2	AMT3	AMT4	AMT5
# of disagree	348	353	376	338	441

Table: Total number of disagreement each AMT has

Performance comparison



(L) AMT 5 was quickly weeded out; eventually settled on the optimal set of AMTs 1, 2, and 4 for most of the time.

(R) CDF of all images' labeling error at the end of this process.

Conclusion

We discussed a quality control problem in labeler market

- How to select the best set of labelers over a sequence of tasks.
 - An algorithm that estimates labeler's quality by comparing against (weighted) majority vote; new regret bound.

Currently under investigation

- Lower bound on the regret in the labeler selection problem.
- Hypothesis testing & coupling argument.

Q & A

Thank you. Any question? http://www.umich.edu/~youngliu